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RPC properties







Differences between RPC and local procedure call

RPC is an extension of the same type of communication to programs running on different computers; single

callee

program

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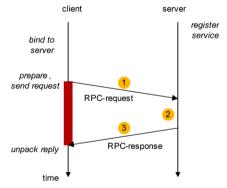
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Neither the client nor the server assume that the procedure call is performed over a network.

Control flow for RPC calls



<u>Differe</u> <u>Ces between RPC and local procedure call</u> <u>Basic RPC characteristics</u>

RPC and OSI

RPC vs message exchange



Definition RPC properties

For an RPC, the caller and the callee run in different processes.

both processes (caller and callee) have

no shared address space.

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single process

answer

thread of execution and transfer of data.

request

procedure

interface between remote

systems

answer

Local vs. remote procedure call

caller

caller

program

no common runtime environment.

different life span of client and server .

Handle errors occurring during a RPC call, e.g. caused by machine crashes or communication failures

RPC-based applications must take communication failures into consideration.

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RPC and OSI



An RPC can be characterized as follows

- 1. uniform call semantics.
- 2. "type-checking" of parameters and results.
- 3. parameter functionality.
- 4. Optimize response times rather than throughput.
- 5. new error cases

bind operation failed; request timed out; arguments are too large

goal is some <u>transparency</u> concerning exception handling and communication failures (relevant for the programmer).

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Integration of the RPC into ISO/OSI protocol stack

layer 7 application layer	client-server model	
layer 6 presentation layer	RPC	hides communication details
layer 5 session layer	message exchange , e.g. request-response protocol	Operating system interface to underlying communication protocols
layer 4 transport layer	transport protocols e.g. TCP/UDP or OSI TP4	transfer of data packets

transport protocols: UDP (User Datagram Protocol) transports data packets without guarantees; TCP (Transmission Control Protocol) verifies correct delivery of data streams.

message exchange: socket interface to the underlying communication protocols.

RPC: hides communication details behind a procedure call and helps bridge heterogeneous platforms.

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RPC vs message exchange





RPC exchange protocols



RPC	message exchange
synchronous (generally)	asynchronous
1 primitive operation (RPC call)	2 primitive operation (send, receive)
messages are configured by RPC system	message specification by programmer

The RPC protocol defines only the structure of the request/answer messages; it does not supply a mechanism for secure data transfer.

several parallel messages possible

RPC exchange protocols

one open RPC

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There are different types of RPC exchange protocols

the request (R) protocol

the request-reply (RR) protocol

the request-reply-acknowledge (RRA) protocol.

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Remote Invocation (RPC/RMI)







Stubs



Issues

Introduction

Distributed applications based on RPC

Remote Method Invocation (RMI)

Servlets

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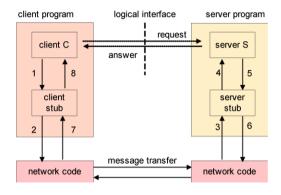
Integration of software handling the communication between components of a distributed application.

Stubs encapsulate the distribution specific aspects.

Stubs represent interfaces.

Client Stub: contains the proxy definition of the remote procedure P.

Server Stub: contains the proxy call for the procedure P.



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Stub functionality









Distributed applications based on RPC



Client and server stubs have the following tasks during client - server interaction.

1. Client stub

specification of the remote service operation; assigning the call to the correct server; representation of the parameters in the transmission format.

decoding the results and propagating them to the client application.

unblocking of the client application.

2. Server stub

decoding the parameter values; determining the address of the service operation (e.g. a table lookup).

invoking the service operation.

prepare the result values in the transmission format and propagate them to the client.

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How to implement distributed applications based on remote procedure calls?



Distributed application

In order to isolate the communication idiosyncrasy of RPCs and to make the network interfaces transparent to the application programmer, so-called stubs are introduced.

Stubs

Stub functionality

Implementing a distributed application

RPC language

Phases of RPC based distributed applications

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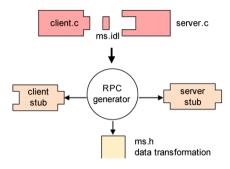




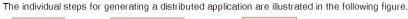
An RPC generator

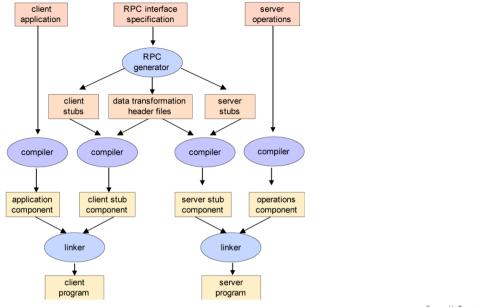
reduces the time necessary for implementation and management of the components of a distributed application.

a declarative interface description is easier to modify and therefore less error-prone.



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Structure of a distributed application



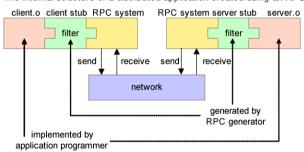




Implementing a distributed application



The internal structure of a distributed application created using an RPC generator is as follows:



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Manual implementation of stubs is error-prone ⇒ use of a *RPC generator* to generate stubs from a declarative specification.

RPC generator

Applying the RPC generator

Structure of a distributed application

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